

OLIMPIADA NATIONALA DE MATEMATICA

ETAPA FINALA - 2016

1.

Fie $n \in \mathbb{N}, n \geq 2$ si numerele reale $a_1, a_2, \dots, a_n \in (0, \infty)$ astfel incat $a_1 \cdot a_2 \cdot \dots \cdot a_n = 1$. Demonstrati ca functia

$$f : (0, \infty) \rightarrow \mathbb{R}, f(x) = (1 + a_1^x)(1 + a_2^x) \dots (1 + a_n^x)$$

este crescatoare.

Solutie.

Efectuand inmultirile se obtine:

$$f(x) = 1 + \sum_{1 \leq i \leq n} a_i^x + \sum_{1 \leq i < j \leq n} a_i^x a_j^x + \sum_{1 \leq i < j < k \leq n} a_i^x a_j^x a_k^x \dots + a_1^x a_2^x \dots a_n^x \quad (1).$$

Deoarece $a_1 \cdot a_2 \cdot \dots \cdot a_n = 1$, expresia functiei f din (1) se poate scrie ca o

suma de expresii de forma $a^x + \frac{1}{a^x}$, unde $a \in (0, \infty)$ (2).

Considerand functia $g : (0, \infty) \rightarrow \mathbb{R}, g(x) = a^x + \frac{1}{a^x}$, pentru $x, t > 0$

avem:

$$g(x+t) - g(x) = \frac{a^{2x+2t} + 1}{a^{x+t}} - \frac{a^{2x} + 1}{a^x} = \frac{a^{2x+2t} - a^{2x+t} + 1 - a^t}{a^{x+t}} = \frac{(a^{2x+t} - 1)(a^t - 1)}{a^{x+t}} \stackrel{x, t > 0}{\geq} 0$$

de unde rezulta g crescatoare (3).

Atunci din (1), (2) si (3) se obtine f crescatoare.

2.

Fie $f : \mathbb{R} \rightarrow \mathbb{R}$ o functie cu proprietatile:

$$(P1) f(x+y) \leq f(x) + f(y),$$

$$(P2) f(tx + (1-t)y) \leq tf(x) + (1-t)f(y),$$

pentru orice $x, y \in \mathbb{R}$ si $t \in [0, 1]$.

a. Demonstrati ca oricare ar fi $a \leq b \leq c \leq d$, astfel incat $d - c = b - a$, are loc inegalitatea $f(b) + f(c) \leq f(a) + f(d)$.

b. Demonstrati ca

$$f(x_1 + \dots + x_{n-1} + x_n) + (n-2)[f(x_1) + \dots + f(x_{n-1}) + f(x_n)] \geq \sum_{1 \leq i < j \leq n} f(x_i + x_j),$$

pentru orice $n \in \mathbb{N}, n \geq 3$ si $x_1, x_2, \dots, x_n \in \mathbb{R}$.

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Solutie.

$$a) \text{Fie } t \in [0,1] \text{ a.i. } b = ta + (1-t)d \text{ si } c = (1-t)a + td \Rightarrow \begin{cases} f(b) \stackrel{ip.}{\leq} tf(a) + (1-t)f(d) \\ f(c) \stackrel{ip.}{\leq} (1-t)f(a) + tf(d) \end{cases}$$

$$\Rightarrow f(b) + f(c) \leq f(a) + f(d).$$

b) Prin inductie folosind urmatoarea schema de rezolvare:

I. $P(3)$ – adevarat;

$$\text{II. } f(x_1 + \dots + x_n + x_{n+1}) + (n-2)[f(x_1) + \dots + f(x_{n-1}) + f(x_n + x_{n+1})] \stackrel{P(n)A}{\geq}$$

$$\sum_{1 \leq i < j \leq n-1} f(x_i + x_j) + \sum_{1 \leq i \leq n-1} f(x_i + x_n + x_{n+1}); \quad \text{III. } \sum_{1 \leq i \leq n-1} f(x_i + x_n + x_{n+1}) \geq$$

$$\sum_{1 \leq i \leq n-1} [f(x_i + x_n) + f(x_i + x_{n+1}) - f(x_i)] + (n-1)[f(x_n + x_{n+1}) - f(x_n) - f(x_{n+1})];$$

IV. $P(n) \Rightarrow P(n+1)$;

Intr-adevar:

I. Fie $x, y, z \in \mathbb{R}$ si $x \geq y \geq 0$ (de exemplu). Avem: $z \leq y + z \leq x + z \leq x + y + z \Rightarrow$ (a)

$$f(y+z) + f(x+z) \stackrel{(a)}{\leq} f(z) + f(x+y+z) \Rightarrow f(x+y) + f(y+z) + f(x+z) \leq$$

$$\stackrel{(P1)}{f(x+y) + f(z) + f(x+y+z)} \leq f(x) + f(y) + f(z) + f(x+y+z);$$

II. Aplicam ipoteza inductiva $P(n)$ - adevarat pentru numerele $x_1, \dots, x_{n-1}, x_n + x_{n+1}$

$$\text{III. } \sum_{1 \leq i \leq n-1} f(x_i + x_n + x_{n+1}) \stackrel{P(3)A}{\geq}$$

$$\sum_{1 \leq i \leq n-1} [f(x_i + x_n) + f(x_i + x_{n+1}) + f(x_n + x_{n+1}) - f(x_i) - f(x_n) - f(x_{n+1})];$$

IV. Din **II** si **III** rezulta:

$$f(x_1 + \dots + x_n + x_{n+1}) + (n-2)[f(x_1) + \dots + f(x_{n-1}) + f(x_n + x_{n+1})] \geq \sum_{1 \leq i \leq n-1} [f(x_i + x_n) + f(x_i + x_{n+1}) - f(x_i)] + (n-1)[f(x_n + x_{n+1}) - f(x_n) - f(x_{n+1})] \Leftrightarrow$$

$$f(x_1 + \dots + x_n + x_{n+1}) + (n-1)[f(x_1) + \dots + f(x_{n-1}) + f(x_n) + f(x_{n+1})] \geq$$

$$\sum_{1 \leq i < j \leq n+1} f(x_i + x_j) \Rightarrow P(n+1)A$$

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3.

- a. In planul complex de origine O , consideram punctele A si B , de afixe nenule a si respectiv b . Aratati ca $S(OAB) = \frac{1}{4} |\bar{a} \cdot b - a \cdot \bar{b}|$, unde $S(OAB)$ reprezinta aria triunghiului OAB .
- b. Fie ABC un triunghi echilateral inscris intr-un cerc C de centru O . Pentru un punct P interior cercului C , notam cu $S(P)$ aria triunghiului avand laturile egale cu distantele de la punctul P la varfurile triunghiului ABC . Fie P_1 si P_2 doua puncte distincte interioare cercului C . Aratati ca $S(P_1) = S(P_2)$ daca si numai daca $OP_1 = OP_2$.

Solutie.

$$S(OAB) = \frac{1}{2} OA \cdot OB \cdot \sin(\angle AOB)$$

$$a) \sin(\arg(z)) = \frac{1}{2 \cdot i} \left(\frac{z - \bar{z}}{|z|} \right) \Rightarrow S(OAB) = \frac{1}{4} |\bar{a}b - a\bar{b}|.$$

$$\sin(\angle AOB) = \left| \sin\left(\arg\frac{a}{b}\right) \right|$$

b) Putem presupune $A(1), B(\varepsilon), C(\varepsilon^2), \varepsilon \in U_3 \setminus \{1\}$.

Fie $P(p), D(p-1), E(\varepsilon(p-\varepsilon)), F(\varepsilon^2(p-\varepsilon^2))$. Folosim urmatoarea schema de rezolvare:

I. $OD = PA, OE = PB, OF = PC;$

II. $S(P) = S(ODE) = S(OEF) = S(OFD) = \frac{\sqrt{3}}{4} (1 - |p|^2);$

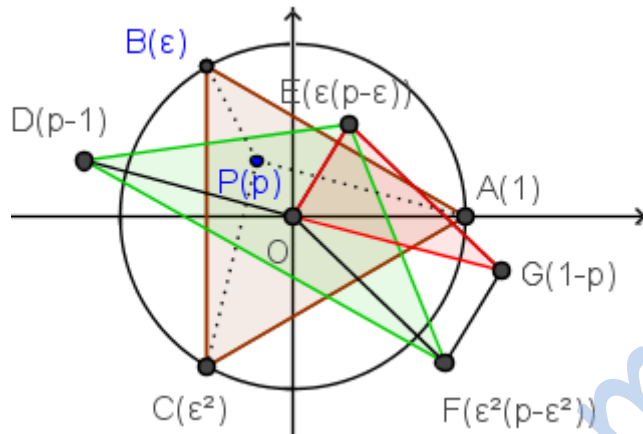
III. $S(P_1) = S(P_2) \Leftrightarrow OP_1 = OP_2.$

Intr-adevar:

I. $OD = |p-1| = |z_P - z_A| = PA, OE = |\varepsilon(p-\varepsilon)| = |p-\varepsilon| = |z_P - z_B| = PB;$

$OF = |\varepsilon^2(p-\varepsilon^2)| = |p-\varepsilon^2| = |z_P - z_C| = PC.$

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II. Fie punctul $G(1-p)$. Avem:

$z_D + z_E + z_F = 0 \Rightarrow O = \text{centrul de greutate al triunghiului } DEF \Rightarrow$

$$S(ODE) = S(OEF) = S(OFD) \stackrel{D(p-1), G(1-p)}{=} S(OEG) \quad (1);$$

$$\left. \begin{array}{l} \text{I} \\ OE = PB, \\ OG = OD = PA, \\ \text{I} \\ \left. \begin{array}{l} OEGF \text{ paralelogram} \\ EG = OF = PC \end{array} \right\} \Rightarrow S(OED) = S(P) \quad (2) \end{array} \right\}$$

$$S(OED) \stackrel{(a)}{=} \frac{1}{4} \left| \overline{\varepsilon(p-\varepsilon)} \cdot (p-1) - \varepsilon(p-\varepsilon)\overline{(p-1)} \right| =$$

$$\frac{1}{4} \left[\varepsilon^2(\bar{p} - \varepsilon^2)(p-1) - \varepsilon(p-\varepsilon)(\bar{p}-1) \right]$$

$$\frac{1}{4} \left| (2\varepsilon+1)(1-|p|^2) \right| = \frac{\sqrt{3}}{4} (1-|p|^2) \stackrel{(1),(2)}{\Rightarrow}$$

$$S(P) = S(ODE) = S(OEF) = S(OFD) = \frac{\sqrt{3}}{4} (1-|p|^2);$$

III. Daca p_1 si p_2 sunt afixele punctelor P_1 respectiv P_2 avem:

$$S(P_1) = S(P_2) \stackrel{\text{II}}{\Leftrightarrow} \frac{\sqrt{3}}{4} (1-|p_1|^2) = \frac{\sqrt{3}}{4} (1-|p_2|^2) \Leftrightarrow |p_1| = |p_2| \Leftrightarrow OP_1 = OP_2.$$

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4.

Oamenii unui trib stravechi foloseau o limba in care cuvintele erau formate doar cu literele A si B . Cercetatorii au descoperit ca pentru oricare doua cuvinte de lungimi egale, exista cel putin trei pozitii corespondente in care literele sunt diferite. De exemplu, cuvintele $ABBAA$ si $AAAAB$ difera in pozitile 2,3 si 5, adica in trei pozitii.

Fie $n \in \mathbb{N}, n \geq 3$. Demonstrati ca in aceasta limba nu pot exista mai mult de $\left\lfloor \frac{2^n}{n+1} \right\rfloor$ cuvinte de lungime n ($[a]$ este partea intreaga a numarului real a)

Solutie.

Facem urmatoarele notatii:

C =multimea tuturor cuvintelor de lungime n care s-ar putea forma in limba respectiva (cuprinde eventual si cuvinte care nu fac parte din limba);

D = multimea tuturor cuvintelor de lungime n din limba respectiva ($D \subset C$);

$d(x, y)$ =numarul de pozitii in care difera doua cuvinte $x, y \in C$;

$C_a = \{x \in C \mid d(x, a) \leq 1\}, a \in C$.

Folosim urmatoarea schema de rezolvare:

I. $card(D) \leq card(C) = 2^n$; **II.** $a, b \in D$ si $a \neq b \Rightarrow C_a \cap C_b = \emptyset$;

III. $card(C_a) = n+1, \forall a \in C$; **IV.** $card(D) \leq \left\lfloor \frac{2^n}{n+1} \right\rfloor$.

Intr-adevar:

I. $D \subset C \Rightarrow card(D) \leq card(C) = 2^n$;

II. Fie prin **R.A.** $a, b \in D$ si $x \in C_a \cap C_b \Rightarrow \begin{cases} d(x, a) \leq 1 \\ d(b, x) \leq 1 \end{cases} \Rightarrow$

$\begin{matrix} a, b \in D \\ d(b, a) \leq 2 \end{matrix} \Rightarrow$ **Contradicție**;

III. C_a contine pe a si inca n cuvinte obtinute prin modificarea lui a in cate o pozitie din cele n ;

$$\bigcup_{a \in D} C_a \subset C \Rightarrow card\left(\bigcup_{a \in D} C_a\right) \stackrel{\text{I}}{\leq} 2^n \Leftrightarrow \sum_{a \in D} card(C_a) \stackrel{\text{II}}{\leq} 2^n \Leftrightarrow card(D) \cdot (n+1) \stackrel{\text{III}}{\leq} 2^n \Leftrightarrow$$

$$card(D) \leq \frac{2^n}{n+1} \quad \Leftrightarrow \quad card(D) \in \mathbb{N} \quad \Leftrightarrow \quad card(D) \leq \left\lfloor \frac{2^n}{n+1} \right\rfloor$$